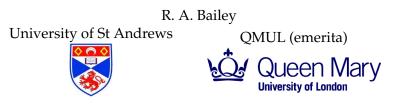
Circular designs with weak neighbour balance



ACCMCC, December 2015

Joint work with Katarzyna Filipiak and Augustyn Markiewicz (Poznan University of Life Sciences), Joachim Kunert (TU Dortmund) and Peter Cameron (St Andrews)

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Wind \rightarrow									
6:0	1	2	3	4	5	6			
5:0	2	4	6	1	3	5			
3:0	4	1	5	2	6	3			
6:0	1	2	3	4	5	6			
5:0	2	4	6	1	3	5			
4:0	3	6	2	5	1	4			
3:0	4	1	5	2	6	3			
2:0	5	3	1	6	4	2			
1:0	6	5	4	3	2	1			

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6:0	1	2	3	4	5	6			
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4:0	3	6	2	5	1	4			
3:0	4	1	5	2	6	3			
2:0	5	3	1	6	4	2			
1:0	6	5	4	3	2	1			

c	•—	#	times	i	is	directly
Sij	.—	uj	pwind	0	f j	

	I	Nin	d -	\rightarrow			_											
6:0	1	2	3	4	5	6	\int_{s} # times <i>i</i> is directly											
5:0	2	4	6	1	3	5	$s_{ij} := $ upwind of j							s_{η} upwind of <i>j</i>				
3:0	4	1	5	2	6	3												
6:0	1	2	3	4	5	6	0 1 2 3 4 5 6											
5:0	2	4	6	1	3	5	$\begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix}$											
4:0	3	6	2	5	1	4												
3:0	4	1	5	2	6	3	$\begin{vmatrix} S = 3 \\ 4 \end{vmatrix} \qquad 0$											
2:0	5	3	1	6	4	2	5 0											
1:0	6	5	4	3	2	1												

	V	Nin	d -	\rightarrow								
6:0	1	2	3	4	5	6	$s_{ij} := $ [#] times <i>i</i> is directly upwind of <i>j</i>					
5:0	2	4	6	1	3	5						
3:0	4	1	5	2	6	3						
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3:0	4	1	5	2	6	3	S = 3 0 0					
2:0	5	3	1	6	4	2	5 0					
1:0	6	5	4	3	2	1	6 \ 0/					

	I	Nin	d -	\rightarrow			
6:0	1	2	3	4	5	6	$\frac{1}{2}$ # times <i>i</i> is directly
5:0	2	4	6	1	3	5	$s_{ij} :=$ upwind of <i>j</i>
3:0	4	1	5	2	6	3	
6:0	1	2	3	4	5	6	0 1 2 3 4 5 0
5:0	2	4	6	1	3	5	$0 \left(\begin{array}{cccccccc} 0 & 2 & 2 & 1 & 2 & 1 \\ 1 & 2 & 2 & 2 & 1 & 2 \end{array} \right)$
4:0	3	6	2	5	1	4	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
3:0	4	1	5	2	6	3	$S = \begin{array}{c ccccccccccccccccccccccccccccccccccc$
2:0	5	3	1	6	4	2	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
1:0	6	5	4	3	2	1	6 2 2 1 2 1 1

1 0

A design with *t* treatments each occurring once in each circular block of size *t* is

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RAB and PJC gave some constructions and non-existence results.

A 0,1-matrix

If we have a design which is weakly neighbour balanced but not neighbour balanced then *S* has zero diagonal, some other entries $\lambda - 1$ and some other entries λ . Put

$$A = S - (\lambda - 1)(J - I).$$

Then

- A is not zero;
- all entries of A are in {0,1};
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- ► $A^{\top}A (\lambda 1)(A + A^{\top})$ is completely symmetric.

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We know something about (some) matrices like this!

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If Type III, then $A^{\top}A$ is not completely symmetric.

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Number the positions in each block 1, 2, ..., starting at the windy end.

Theorem

If a WNBD has the property that each numbered position has all treatments equally often, then it either is a NBD or has Type I.

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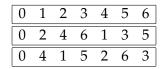
$$0 \neq y^2 \in Z_t \qquad x+1$$

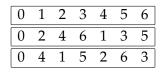
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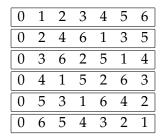
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 $t = 3 \checkmark$, but too small to separate direct effects from upwind effects

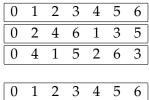
 $t = 7 \checkmark$, see next slide

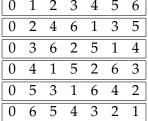






 $t = 11 \checkmark$





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t = 15? RAB tried using *A* as the incidence matrix of PG(3, 2) and proved that it is impossible.

Reid and Brown give the following doubling construction.

$$A_2 = \left(egin{array}{ccc} A_1^{ op} & 0_t & A_1 + I_t \ 1_t^{ op} & 0 & 0_t^{ op} \ A_1 & 1_t & A_1 \end{array}
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If A_1 is Type I for *t* then A_2 is Type I for 2t + 1.

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Doing this with t = 7 gives a doubly regular tournament Γ_2 on 15 vertices with an automorphism π of order 7. If we can find a Hamiltonian cycle φ which has no edge in common with any of $\pi^i(\varphi)$ for i = 1, ..., 6, then $\varphi, \pi(\varphi), ..., \pi^6(\varphi)$ make a WNBD.

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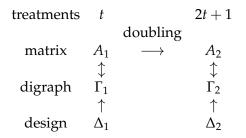
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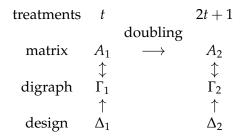
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RAB tried and failed to do this by hand. PJC used GAP, and found 120 solutions. KF put this A_2 into Mathematica and asked it to find Hamiltonian decompositions.





Could we go directly from Δ_1 to Δ_2 ?

Type I designs with rows and columns

Suppose that $t \equiv 3 \mod 4$ and t is a prime power. Let x be a primitive element of GF(t). In the circular sequence

$$(1, x, x^2, x^3, \dots, x^{t-1})$$

the successive differences give all non-zero elements of GF(t).

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Put
$$\phi = (x, 1, 0, x^2, x^3, \dots, x^{t-1}).$$

If $t \neq 3$ then the number of non-zero squares in the successive differences of ϕ is one different from the number of non-squares in the successive differences of ϕ .

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The t(t-1)/2 sequences $s\phi + i$ where *s* is a non-zero square in GF(*t*) and $i \in GF(t)$ give a weakly neighbour-balanced design in which every treatment occurs (t-1)/2 times in each numbered position.

Type II: $A^{\top}A$ is completely symmetric and $\lambda = 1$

Now we can regard *A* as the incidence matrix of a 2-design, with blocks labelled so that the diagonal is zero.

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We can also take advantage of symmetry to find a single Hamiltonian cycle whose images under a group of automorphisms of Γ give the blocks of the WNBD. Now we can regard *A* as the incidence matrix of a 2-design, with blocks labelled so that the diagonal is zero.

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If *A* is itself symmetric then it is the adjacency matrix of a strongly regular graph in which every pair of distinct vertices have the same number of common neighbours (for example, the Shrikandhe graph and the Clebsch graph).

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Again, familiar tricks and use of symmetry give us WNBDs.

Type III: $A^{\top}A - (\lambda - 1)(A + A^{\top})$ is completely symmetric, but $A^{\top}A$ and $(A + A^{\top})$ are not

If A_1 has Type I for *t* treatments then

$$\begin{pmatrix} A_{1} & A_{1} + I_{t} & \dots & A_{1} + I_{t} \\ A_{1} + I_{t} & A_{1} & \dots & A_{1} + I_{t} \\ \vdots & \vdots & \ddots & \vdots \\ A_{1} + I_{t} & A_{1} + I_{t} & \dots & A_{1} \end{pmatrix}$$

and
$$\begin{pmatrix} 0 & 1_{t}^{\top} & 0 & 0_{t}^{\top} \\ 0_{t} & A_{1} & 1_{t} & A_{1}^{\top} \\ 0 & 0_{t}^{\top} & 0 & 1_{t}^{\top} \\ 1_{t} & A_{1}^{\top} & 0_{t} & A_{1} \end{pmatrix}$$
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has Type III for *mt* treatments with $\lambda = m(t+1)/4$

has Type III for 2(t + 1) treatments with $\lambda = (t + 1)/2$.

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The second type is the adjacency matrix of what Cameron and Babai call an S-digraph.

Type III: $A^{\top}A - (\lambda - 1)(A + A^{\top})$ is completely symmetric, but $A^{\top}A$ and $(A + A^{\top})$ are not

If A_1 has Type I for *t* treatments then

$$\begin{pmatrix} A_1 & A_1 + I_t & \dots & A_1 + I_t \\ A_1 + I_t & A_1 & \dots & A_1 + I_t \\ \vdots & \vdots & \ddots & \vdots \\ A_1 + I_t & A_1 + I_t & \dots & A_1 \end{pmatrix} \text{ has Type III for } mt \text{ treatments} \\ \text{with } \lambda = m(t+1)/4 \\ \text{and} \quad \begin{pmatrix} 0 & 1_t^\top & 0 & 0_t^\top \\ 0_t & A_1 & 1_t & A_1^\top \\ 0 & 0_t^\top & 0 & 1_t^\top \\ 1_t & A_1^\top & 0_t & A_1 \end{pmatrix} \text{ has Type III for } 2(t+1) \text{ treatments} \\ \text{with } \lambda = (t+1)/2. \end{cases}$$

The second type is the adjacency matrix of what Cameron and Babai call an S-digraph. t = 3 leads to the only Type III WNBDs (t = 6 and t = 8) found by KF and AM.

Type III doubling (or multiplying) constructions

Again, is there a way of going directly from the smaller design to the larger one?