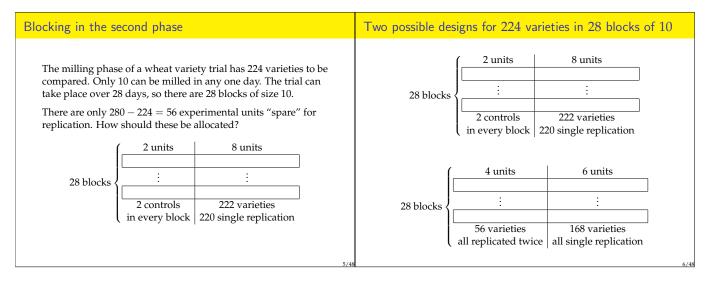


In breeding trials of new varieties, typically there is very little	"on any given field agricultural operations, at least for
seed of each of the new varieties. Traditionally, an experiment has one plot for each new variety and several plots for a well-established "control": for example, 30 new varieties on one plot each and one control on 8 plots. In the last 10 years, Cullis and colleagues in Australia (and independently Bueno and Gilmour) have suggested replacing many occurrences of the the control by double replicates of a small number of new varieties: for example, 24 new varieties with one plot each, 6 new varieties with two plots each, and the control on two further plots. This is an improvement if there are no blocks.	 Centuries, have followed one of two directions, which are usually those of the rows and columns; consequently streaks of fertility, weed infestation, etc., do, in fact, occur predominantly in those two directions." R. A. Fisher, letter to H. Jeffreys, 30 May 1938 (selected correspondence edited by J. H. Bennett) (This assumption is dubious for field trials in Australia.) If field operations have been primarily in one direction for a long time, then it is reasonable to divide the fields into blocks whose length runs along that direction.



The problem Linear model, estimation and variance We measure the response *Y* on each unit in each block. If that unit has variety *i* and block *D*, then we assume that $Y = \tau_i + \beta_D +$ random noise, We are given *b* blocks of size *k*. We are given *v* varieties. Assume that where the random noise is independently normally distributed average replication $= \bar{r} = \frac{bk}{v} \le 2.$ with zero mean and constant variance σ^2 . We want to estimate all the simple differences $\tau_i - \tau_i$. How should we allocate varieties to blocks? What makes a block design good? Put variance of the best linear unbiased estimator $V_{ij} \sigma^2$ = for $\tau_i - \tau_j$. We want all the V_{ij} to be small. Optimality An example with 5n + 10 varieties in 5 blocks of size 4 + n

Apart from the constant multiple σ^2 ,

$$V_{ii}$$
 = variance of the BLUE for $\tau_i - \tau_i$.

Put

$$V_T = \sum_{i=1}^{v-1} \sum_{j=i+1}^{v} V_{ij}$$
 = sum of variances of variety differences.

Definition

For given values of b (the number of blocks), k (the size of the blocks) and v (the number of varieties), a block design is A-optimal if it minimizes V_T .

9	0	1	2	E_1	• • •	E_n

 $3 \quad 4 \quad 5 \quad 6 \quad B_1 \quad \cdots \quad B_n$

 $1 \ 2 \ 3 \ 4 \ A_1 \ \cdots$

 $5 \ 6 \ 7 \ 8 \ C_1 \ \cdots$

 $7 8 9 0 D_1 \cdots$

 A_n

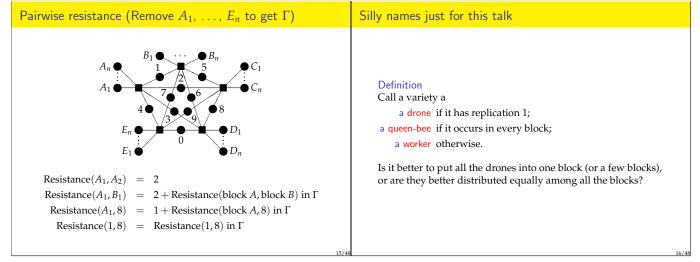
 C_n

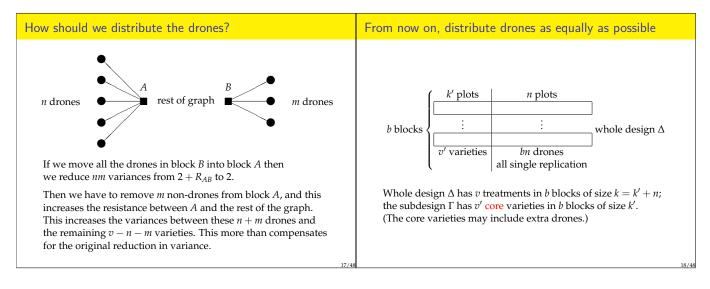
 D_n

How do we calculate pairwise variances in a generic design?

Levi graph Levi graph: example $2 \ 3 \ 4 \ A_1 \ \cdots$ A_n $B_1 \cdots$ B_n 5 6 7 5 8 9 C₁ C_n 2 . . . The Levi graph of the block design has 3 6 8 0 D₁ ... D_n one vertex for each variety 9 0 E₁ E_n 7 . . . one vertex for each block one edge for each plot (aka experimental unit), so that the edge for plot ω joins the vertex for the variety on ω to the vertex for the block containing ω . D. 12/48

We can consider the Levi graph as an electrical network with a 1-ohm resistance in each edge. Connect a 1-volt battery between vertices <i>i</i> and <i>j</i> . Current flows in the network, according to these rules. 1. Ohm's Law: In every edge, voltage drop = current × resistance = current. 2. Kirchhoff's Voltage Law: The total voltage drop from one vertex to any other vertex is the same whichever path we take from one to the other. 3. Kirchhoff's Current Law: At every vertex which is not connected to the battery, the total current coming in is equal to the total current going out. Find the total current <i>I</i> from <i>i</i> to <i>j</i> , then use Ohm's Law to define the effective resistance R_{ij} between <i>i</i> and <i>j</i> as 1/ <i>I</i> .	Electrical networks	Electrical networks: variance
define the effective resistance R_{ij} between <i>i</i> and <i>j</i> as $1/I$. matrix inversion if the graph is sparse.	 with a 1-ohm resistance in each edge. Connect a 1-volt battery between vertices <i>i</i> and <i>j</i>. Current flows in the network, according to these rules. 1. Ohm's Law: In every edge, voltage drop = current × resistance = current. 2. Kirchhoff's Voltage Law: The total voltage drop from one vertex to any other vertex is the same whichever path we take from one to the other. 3. Kirchhoff's Current Law: At every vertex which is not connected to the battery, the total current coming in is equal to the total current going out. 	TheoremIf R_{ij} is the effective resistance between variety vertices i and j in the Levi graph then $R_{ij} = V_{ij}$.Put: V_{CD} = variance of BLUE of $\beta_C - \beta_D$ for blocks C and D, V_{iC} = variance of BLUE of $\tau_i + \beta_C$ for variety i and block C.TheoremIf R_{CD} and R_{iC} are the effective resistances between vertices C and D, and between i and C respectively, in the Levi graph then $R_{CD} = V_{CD}$ and $R_{iC} = V_{iC}$.
10/10		matrix inversion if the graph is sparse.





Sum of the pairwise variances

Sum of variances in whole design if Γ is equi-replicate

Theorem (cf. Herzberg and Jarrett, 2007) *If there are n drones in each block of* Δ *,* and the core design Γ has v' varieties in b blocks of size k'

then the sum of the variances of variety differences in Δ

 $= V_T(\Delta) = bn(bn + v' - 1) + V_T + nV_{BT} + n^2V_B,$ where

$$V_{T}$$
 = the sum of the variances of variety differences in Γ

 $V_B = the sum of the variances of block differences in \Gamma$

 V_{BT} = the sum of the variances of sums of one treatment and one block in Γ .

$V_T(\Delta) = bn(bn + v' - 1) + V_T + nV_{BT} + n^2V_B$

 V_T = the sum of the variances of variety differences in Γ

 V_B = the sum of the variances of block differences in Γ

 V_{BT} = the sum of the variances of sums of one treatment and one block in Γ .

If Γ is equi-replicate with replication r' then

$$\begin{array}{lll} \frac{k'}{b}V_B - b &=& \frac{r'}{v'}V_T - v';\\ V_{BT} &=& \frac{2b}{v'}V_T + \frac{v'}{k'}(b - v' - 1), \end{array}$$

and so V_B and V_{BT} are both increasing functions of V_T .

Consequence

For a given choice of k', use the core design Γ which minimizes V_T .

Sum of variances in whole design if there are many drones	Strategy
$V_T(\Delta) = bn(bn + v' - 1) + V_T + nV_{BT} + n^2V_B$ $V_T = \text{the sum of the variances of variety differences in } \Gamma$ $V_B = \text{the sum of the variances of block differences in } \Gamma$ $V_{BT} = \text{the sum of the variances of sums of one treatment and one block in } \Gamma.$ Consequence If n is large, we need to focus on reducing V_B,	Given <i>b</i> , <i>v</i> and <i>k</i> , how do we find an A-optimal design for <i>v</i> varieties in <i>b</i> blocks of size <i>k</i> when $\frac{bk}{2} \le v \le b(k-1) + 1?$ Average replication ≤ 2 Maximum <i>v</i> for estimability Case 1. <i>b</i> = 2 or <i>b</i> = 3 (very small <i>b</i>). Case 2. <i>v</i> = <i>b</i> (<i>k</i> - 1) + 1 (very large <i>v</i>). Case 3. <i>v</i> = <i>b</i> (<i>k</i> - 1) (very large <i>v</i>).
so it may be best to increase the number of drones and decrease k' (the size of blocks in the core design Γ), so that average replication within Γ is more than 2.	Case 4. $k < b - 1$ (small k). Case 5. $k \ge b - 1$.

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Case 1. Only 2 blocks, of size k	Case 1 continued. 3 blocks of size k
Morgan and Jin (2007) showed that the A-optimal designs are those with 2 <i>n</i> drones and <i>q</i> queen bees, where $q = 2k - v$ and $n = k - q$. $ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	Using exhaustive (and exhausting (and tedious)) case-by-case analysis, RAB has shown that the A-optimal designs are as follows when <i>v</i> is divisible by 3 (and presumably small changes deal with the other cases). There are $3w$ workers and $3n$ drones, where $3w = 3k - v$ and $n = k - 2w$.
They also showed that, when $2k - v$ is comparatively large, the MV-optimal designs (those designs that minimize the maximum of the pairwise variances V_{ij}) have all the drones in the same block.	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$

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Case 2. v = b(k-1) + 1

This is the maximum number of varieties that can be tested in \boldsymbol{b} blocks of size \boldsymbol{k} with all comparisons estimable.

Mandal, Shah and Sinha (1991), for k = 2, and Bailey and Cameron (2013), for general block size, showed that, no matter how many blocks there are, the A-optimal design has the following form.

1	A_1	A_2	A_3		A_{k-1}
1	B_1	B ₂	B ₃		B_{k-1}
1	C_1	<i>C</i> ₂	<i>C</i> ₃		C_{k-1}
1	D_1	D ₂	D ₃		D_{k-1}
1	E_1	<i>E</i> ₂	<i>E</i> ₃		E_{k-1}
1 queen		<i>v</i> -	- 1 dı	rones	;

Case 3. v = b(k-1)

The A-optimal design by Krafft and Schaefe small <i>k</i> and <i>b</i>	er (1997).	r all cases then increase b
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c c c c c c c c c c c c c c c c c c c $

Youden and Connor (1953) had recommended chain designs.

Case 3. $v = b(k-1)$ revisited	Case 4. $k < b-1$
Theorem Consider a design with b blocks of size 2. For $2 \le s \le b$, let Γ_s be the design consisting of a chain of length s, one of whose varieties is in all blocks outside the chain, while all other varieties are drones. Then $V_B(\Gamma_s) = \frac{1}{6}[-s^3 + 2bs^2 - (6b - 4)s + 6b^2 - 5b].$	For various values of $k_i \leq k$, find the best core design Γ_i for v'_i varieties in <i>b</i> blocks of size k_i . (For equi-replicate core designs, it is often easier to find the best dual design, which is obtained by interchanging the roles of blocks and varieties.) $V_T(\Gamma_i) =$ the sum of the variances of variety differences in Γ_i $V_B(\Gamma_i) =$ the sum of the variances of block differences in Γ_i
Consequence 1. If $b = 3$ then $V_B(\Gamma_2) > V_B(\Gamma_3)$ so there is no need for queens.	$V_{BT}(\Gamma_i)$ = the sum of the variances of sums of one treatment and one block in Γ_i .
2. If $b = 4$ then $V_B(\Gamma_2) = V_B(\Gamma_4) < V_B(\Gamma_3)$, but $V_T(\Gamma_2) > V_T(\Gamma_4)$ and $V_{BT}(\Gamma_2) > V_{BT}(\Gamma_4)$, so do not use Γ_2 or Γ_3 (no need for queens).	If there are n_i drones in each block then, in the whole design Δ , $V_T(\Delta) = bn_i(bn_i + v'_i - 1) + V_T(\Gamma_i) + n_i V_{BT}(\Gamma_i) + n_i^2 V_B(\Gamma_i)$
3. If $b \ge 5$ then $V_B(\Gamma_2) < V_B(\Gamma_3) < \cdots < V_B(\Gamma_b)$, so we need to use smaller chains as v gets larger.	Use this formula to find the core design with the smallest $V_T(\Delta)$.

Case 4 continued. $k < b - 1$	Case 4 cor	ntinued. $k =$	$= 4 < b - 1, V_l$	$_{3} \div b(b -$	1)/2
Case 4 continued. $k < b - 1$ If there are n_i drones in each block then, in the whole design Δ , $V_T(\Delta) = bn_i(bn_i + v'_i - 1) + V_T(\Gamma_i) + n_i V_{BT}(\Gamma_i) + n_i^2 V_B(\Gamma_i)$ As the number of varieties increases, it becomes more important to choose Γ_i with a small value of $V_B(\Gamma_i)$.	Case 4 con k_i b = 6 b = 7 b = 8 b = 9 b = 10 b = 11 b = 12 b = 13 b = 14	Best Γ ₁ 2 2 queens,	$= 4 < b - 1, V_{1}$ a design for <i>b</i> blocks Γ_{2} 3 2 queens, 2 workers (rep 2) 1^{-}		
25/4	If $b \ge 14$		1^- omes better than Γ_4 creases, Γ_1 and Γ_2 b		1.08 r than Γ_3 .

Case 4 continued. $k < b-1$ when $b=8$	Case 4 continued. $k = 5$ and $k = 6$ when $b =$	= 8
$k = 6, \text{ and } 24 \text{ varieties, all workers, all replicated twice.}$ $\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	

Case 4 continued. $k = 5$ and k = 5 24 varieties: 16 workers, 8 drones 1 2 3 4 A_1 5 6 7 8 B_1 9 10 11 12 C_1 13 14 15 16 D_1	k = 6 when b = 8 k = 6 32 varieties: 8 workers, 24 drones $1 2 4 A_1 A_2 A_3$ $2 3 5 B_1 B_2 B_3$ $3 4 6 C_1 C_2 C_3$ $4 5 7 D_1 D_2 D_3$	Case 4 continued. $k = 5$ and k = 5 28 varieties: 12 workers, 16 drones $1 \ 2 \ 3 \ A_1 \ A_2$ $1 \ 4 \ 5 \ B_1 \ B_2$ $4 \ 6 \ 7 \ C_1 \ C_2$	d $k = 6$ when $b = 8$ k = 6 36 varieties: 12 workers, 24 drones $1 2 3 A_1 A_2 A_3$ $1 4 5 B_1 B_2 B_3$ $4 6 7 C_1 C_2 C_3$
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$

Now we are also in Case 3 aga	ain!	
k = 5 32 varieties: 1 queen, 1 worker, 30 drones $1 2 A_1 A_2 A_3$ $1 2 B_1 B_2 B_3$ $1 C_1 C_2 C_3 C_4$ $1 D_1 D_2 D_3 D_4$ $1 E_1 E_2 E_3 E_4$ $1 F_1 F_2 F_3 F_4$ $1 G_1 G_2 G_3 G_4$ $1 H_1 H_2 H_3 H_4$	k = 6 40 varieties: 1 queen, 1 worker, 38 drones $1 2 A_1 A_2 A_3 A_4$ $1 2 B_1 B_2 B_3 B_4$ $1 C_1 C_2 C_3 C_4 C_5$ $1 D_1 D_2 D_3 D_4 D_5$ $1 E_1 E_2 E_3 E_4 E_5$ $1 E_1 E_2 F_3 F_4 F_5$ $1 G_1 G_2 G_3 G_4 G_5$ $1 H_1 H_2 H_3 H_4 H_5$	Now $v \ge \frac{bk}{2} \ge \frac{b(b-1)}{2}$. For simplicity, assume that <i>b</i> divides 2 <i>v</i> , and put $n = \frac{2v - bk}{b}$. Then <i>n</i> is the minimum number of drones per block. Let Γ_0 be the design for $b(b-1)/2$ varieties replicated twice in <i>b</i> blocks of size $b - 1$ in such a way that there is one variety in common to each pair of blocks. This is A-optimal for these numbers.

Case 5 continued. $k \geq b-1$ and $v \geq b(b-1)/2$	Case 5. Example: $b=8$ and $k=15$ (so $60 \le v \le 113$)
n = minimal number of drones per block.	60 varieties: all workers
If $k - n \ge b - 1$ then	
1. put n drones in each block;	1 2 3 4 5 6 7 29 30 31 32 33 34 35 57
 <i>put in one copy of</i> Γ₀; <i>put in as many further copies of</i> Γ₀ <i>as possible;</i> 	1 8 9 10 11 12 13 29 36 37 38 39 40 41 57
4. in any remaining space,	2 8 14 15 16 17 18 30 36 42 43 44 45 46 58
use a good design for workers with replication 2 (so long as there is at least one copy of Γ_0 ,	3 9 14 19 20 21 22 31 37 42 47 48 49 50 58
it probably doesn't make much difference which	4 10 15 19 23 24 25 32 38 43 47 51 52 53 59
one is used).	5 11 16 20 23 26 27 33 39 44 48 51 54 55 59
Otherwise we are back in the same situation as Case 4 ($k < b - 1$),	6 12 17 21 24 26 28 34 40 45 49 52 54 56 60
but there are more drones necessary so it is more likely that we will have to move towards a	7 13 18 22 25 27 28 35 41 46 50 53 55 56 60
core design with replication 3 or more, or even queens.	one copy of Γ_0 another copy of Γ_0

Case 5. Example: $b = 8$ and $k = 15$ (so $60 \le v \le 113$)	Case 5. Example: $b=8$ and $k=15$ (so $60 \le v \le 113$)
76 varieties: 44 workers, 32 drones	92 varieties: 28 workers, 64 drones
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$
39/48	40/48

Case 5. Example: $b=8$ and $k=15$ (so $60 \le v \le 113$)	Case 5. Example: $b=8$ and $k=15$ (so $60 \le v \le 113)$
104 varieties: 8 workers, 96 drones	108 varieties: 4 workers, 104 drones
(not 16 workers and 88 drones;	(not 12 workers and 96 drones;
this is the phase-change that we saw before with $k = 6$)	the larger block size forces us past another phase change)
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$	$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$
workers drones	workers drones
rep. 3	rep. 4

Health Warning	References: more even replication
The overall message is that there can be phase changes as the spare capacity for replication $(bk - v)$ decreases. Therefore it is necessary to compare core designs Γ_i with different block size k_i . Although this overall message is correct, no one has checked the arithmetic in the examples presented, so individual cases may be wrong.	 J. S. S. Bueno Filho and S. G. Gilmour: Planning incomplete block experiments when treatments are genetically related. <i>Biometrics</i>, 59, (2003), 375–381. B. R. Cullis, A. B. Smith and N. E. Coombes: On the design of early generation variety trials with correlated data. <i>Journal of Agricultural, Biological and Environmental Statistics</i>, 11, (2006), 381–393. A. B. Smith, P. Lim and B. R. Cullis: The design and analysis of multi-phase plant breeding experiments. <i>Journal of Agricultural Science</i>, 144, (2006), 393–409.

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